

## Research Article

# An Algorithm to Compute the H-Bases for Ideals of Subalgebras

Rabia,<sup>1</sup> Muhammad Ahsan Binyamin ,<sup>1</sup> Nazia Jabeen,<sup>2</sup> Adnan Aslam ,<sup>3</sup>  
and Kraidi Anoh Yannick <sup>4</sup>

<sup>1</sup>Department of Mathematics, GC University, Faisalabad, Pakistan

<sup>2</sup>Department of Mathematical Sciences, Institute of Business Administration, Karachi, Pakistan

<sup>3</sup>Department of Natural Sciences, University of Engineering and Technology (RCET), Lahore, Pakistan

<sup>4</sup>UFR of Mathematics and Informatics, University Félix Houphouët Boigny, Abidjan, Côte d'Ivoire

Correspondence should be addressed to Kraidi Anoh Yannick; kayanoh2000@yahoo.fr

Received 24 April 2021; Accepted 16 June 2021; Published 7 July 2021

Academic Editor: Qamar Din

Copyright © 2021 Rabia et al. This is an open access article distributed under the Creative Commons Attribution License, which permits unrestricted use, distribution, and reproduction in any medium, provided the original work is properly cited.

The concept of H-bases, introduced long ago by Macaulay, has become an important ingredient for the treatment of various problems in computational algebra. The concept of H-bases is for ideals in polynomial rings, which allows an investigation of multivariate polynomial spaces degree by degree. Similarly, we have the analogue of H-bases for subalgebras, termed as SH-bases. In this paper, we present an analogue of H-bases for finitely generated ideals in a given subalgebra of a polynomial ring, and we call them “HSG-bases.” We present their connection to the SAGBI-Gröbner basis concept, characterize HSG-basis, and show how to construct them.

## 1. Introduction

The concept of H-bases, introduced long ago by Macaulay [1], is based solely on homogeneous terms of a polynomial. In [2], an extension of Buchberger’s algorithm is presented to construct H-bases algorithmically. Some applications of H-bases are given in [3]; in addition, many of the problems in applications which can be solved by the Gröbner technique can also be treated successfully with H-bases. The concept of H-basis for ideals of a polynomial ring over a field  $K$  can be adopted in a natural way to  $K$ -subalgebras of a polynomial ring. In [4], SH-basis (Subalgebra Analogue to H-basis for Ideals) for the  $K$ -subalgebra of  $K[x_1, \dots, x_n]$  is defined. The properties of SH-bases are typically similar to H-basis results [3]. Like H-bases, the concept of SH-basis is also tied to homogeneous polynomials. In this paper, we will present an analogue to H-bases for ideals in a given subalgebra of a polynomial ring, and we call them “HSG-bases.”

The paper is organized as follows. In Section 2, we briefly describe the underlying concept of grading which leads to SAGBI-Gröbner bases and HSG-basis. Then, we give the notion of  $si$ -reduction, which is one of the key ingredients for the characterization and construction of HSG-basis.

After setting up the necessary notation, we present the  $si$ -reduction algorithm (see Algorithm 1). Also, here we present some properties characterizing HSG-basis (Theorem 1). In Section 3, we present a criterion through which we can check that the given system of polynomials is an HSG-basis of the subalgebra it generates (Theorem 2), and further on the basis of this theorem, we present an algorithm for the construction of HSG-basis (Algorithm 2).

## 2. HSG-Bases and SAGBI-Gröbner Bases

Here and in the following sections we consider polynomials in  $n$  variables  $x_1, \dots, x_n$  with coefficients from a field  $K$ . For short, we write

$$P := K[x_1, \dots, x_n]. \quad (1)$$

If  $G$  is a subset of subalgebra  $\mathcal{A}$  in  $K[x_1, \dots, x_n]$ , then the set

$$I := \left\{ \sum_{g \in G} h_g g \mid h_g \in \mathcal{A} \text{ and only finitely many } h_g \neq 0 \right\} \quad (2)$$

**Input:** a subalgebra  $\mathcal{A}$ , a finite subset  $G \subset \mathcal{A}$ , and a polynomial  $f \in \mathcal{A}$ .

**Output:** a polynomial  $h$  such that  $f \xrightarrow{G_{\mathcal{A}}^*} h$ .

- (1)  $h := f$ .
- (2) While ( $h \neq 0$  and  $G_h = \{\sum_i a_i g_i \mid M^{(H)}(\sum_i a_i g_i) = M^{(H)}(h) \neq \emptyset\}$ ; where  $a_i \in \mathcal{A}$  and  $g_i \in G$ ).
- (3) Choose  $\sum_i a_i g_i \in G_h$ .
- (4)  $h := h - \sum_i a_i g_i$  and continue at 2.

ALGORITHM 1: Algorithm to compute si-reduction

**Input:** a subalgebra  $\mathcal{A}$  and a finite subset  $G \subset \mathcal{A}$ .

**Output:** HSG-basis  $H$  for  $\langle G \rangle_{\mathcal{A}}$ .

- (1)  $H := G$ , Old( $H$ ):  $= \emptyset$ .
- (2)  $H = \{h_1, \dots, h_s\}$ .
- (3) While ( $H \neq \text{Old}(H)$ ) do
- (4) Compute  $Q$ , an  $M^{(H)}$ -generating set for  $\text{syz}(M^{(H)}(H))$ .
- (5) Compute  $P := \{\sum_{i=1}^s q_i h_i \mid (q_i)_{i=1}^s \in Q\}$ .
- (6) Compute  $\text{red}(P) := \{\text{final si-reduction via } H \text{ of every element of } P\} - \{0\}$ .
- (7) Old( $H$ ):  $= H \cup \text{red}(P)$ .

ALGORITHM 2: Algorithm for the construction of HSG basis.

is the ideal of  $A$  in  $P$  generated by  $G$  and we write it shortly as  $\langle G \rangle_{\mathcal{A}}$ . In this section, we want to introduce HSG-bases and discuss some of their properties. This concept is very similar to the concept of SAGBI-Gröbner bases. Therefore, we will briefly explain the underlying common structure. Let  $\Gamma$  denote an ordered monoid, i.e., an abelian semigroup under an operation  $+$ , equipped with a total ordering  $>$  such that, for all  $\alpha, \beta, \gamma \in \Gamma$ ,

$$\alpha > \beta \Rightarrow \alpha + \gamma > \beta + \gamma. \quad (3)$$

A direct sum,

$$P = \bigoplus_{\gamma \in \Gamma} P_{\gamma}^{(\Gamma)}, \quad (4)$$

is called grading (induced by  $\Gamma$ ) or briefly a  $\Gamma$ -grading if for all  $\alpha, \beta \in \Gamma$ ,

$$f \in P_{\alpha}^{(\Gamma)}, g \in P_{\beta}^{(\Gamma)} \Rightarrow f \cdot g \in P_{\alpha+\beta}^{(\Gamma)}. \quad (5)$$

Since the decomposition above is a direct sum, each polynomial  $f \neq 0$  has a unique representation.

$$f = \sum_{i=1}^s f_{\gamma_i}, \quad 0 \neq f_{\gamma_i} \in P_{\gamma_i}^{(\Gamma)}. \quad (6)$$

Assuming that  $\gamma_1 > \gamma_2 > \dots > \gamma_s$ , the  $\Gamma$ -homogeneous term  $f_{\gamma_1}$  is called the maximal part of  $f$ , denoted by  $M^{(\Gamma)}(f) := f_{\gamma_1}$ , and  $f - M^{(\Gamma)}(f)$  is called the  $d$ -reductum of  $f$ . For  $G \subset \mathcal{A}$ ,  $M^{(\Gamma)}(G) := \{M^{(\Gamma)}(g) \mid g \in G\}$ .

There are two major examples of gradings. The first one is grading by degrees:

$$P_d^{(\Gamma)} = \{p \in P \mid p \text{ is homogeneous of degree } d\}, \quad \forall d \in \mathbb{N}. \quad (7)$$

Here,  $\Gamma = \mathbb{N}$  with the natural total ordering. This grading is called the  $H$ -grading because of the homogeneous polynomials. Therefore, we also write  $H$  in place of this  $\Gamma$ . The space of all polynomials of degree at most  $d$  can now be written as

$$P_d := \bigoplus_{k=0}^d P_k^{(H)}. \quad (8)$$

The maximal part of a polynomial  $f \neq 0$  is its homogeneous form of highest degree,  $M^{(H)}(f)$ . For simplicity, let  $M^{(H)}(0) := 0$ .

**Definition 1.** A subset  $G = \{g_1, \dots, g_s\} \subset \mathcal{A}$  (subalgebra) is called HSG-basis for the ideal  $I_{\mathcal{A}} \subset \mathcal{A}$ , if for all  $0 \neq f \in I_{\mathcal{A}}$ ,

$$\begin{aligned} \exists h_1, \dots, h_s \in \mathcal{A}: f &= \sum_{i=1}^s h_i g_i, \deg(f) = \max_{i=1}^s \{\deg(h_i g_i)\} \text{ (Note that this condition is not obvious, } -x^3 y^3 + x^4 \\ &= (x^2)(x^3 y + x^2) + (-xy)(x^4 + x^2 y^2) \text{ see in } K[x^2, xy]). \end{aligned} \quad (9)$$

The representation for  $f$  in (9) is also called its HSG representation with respect to  $G$ .

Note that HSG-basis for ideal in a subalgebra is also a generating set of it. To obtain more insights into HSG-bases,

we will give some equivalent definitions. First, we need a more technical notion.

**Definition 2.** For given  $f, f_1, \dots, f_m$ , we say that  $f$  si-reduces to  $\tilde{f}$  with respect to  $F = \{f_1, \dots, f_m\}$  in  $\mathcal{A}$  if

$$\tilde{f} = f - \sum_{i=1}^m a_i f_i, \deg(\tilde{f}) < \deg(f) \quad (10)$$

holds with polynomials  $a_i \in \mathcal{A}$  satisfying  $\deg(a_i f_i) \leq \deg(f)$ . We write it as  $f \xrightarrow{F_{\mathcal{A}}} \tilde{f}$ . By  $\xrightarrow{F_{\mathcal{A}}, *}$  we denote the transitive closure of the binary relation  $\xrightarrow{F_{\mathcal{A}}}$ .

The concept of *si*-reduction plays an important role in the characterization and construction of HSG-basis. For  $f \in \mathcal{A}$  and  $G \subset \mathcal{A}$ , the following algorithm computes  $h$  such that  $f \xrightarrow{G_{\mathcal{A}}, *} h$  (i.e.,  $f$  reduces to  $h$  completely).

We note that such an element  $a_i$  in the subalgebra  $\mathcal{A}$  can easily be determined as in the case of reduction in polynomial ring. We also note that  $\deg(h - \sum_i a_i g_i)$  is strictly smaller than the  $\deg(h)$  (by the choice of  $\sum_i a_i g_i$ ). This shows that Algorithm 1 always terminates.

**Theorem 1.** *Let  $G = \{g_1, \dots, g_s\} \subset \mathcal{A}$  (subset of subalgebra  $\mathcal{A}$ ) and  $I_{\mathcal{A}}$  be an ideal of  $\mathcal{A}$ . Then, the following conditions are equivalent:*

- (1)  $G$  is an HSG-basis for the ideal  $I_{\mathcal{A}}$ .
- (2)  $\langle M^{(H)}(g_1), \dots, M^{(H)}(g_s) \rangle_{K[M^{(H)}]} = \langle M^{(H)}(f) | f \in I_{\mathcal{A}} \rangle_{K[M^{(H)}]}$ .
- (3) For all  $f \in I$ ,  $f \xrightarrow{G_{\mathcal{A}}, *} 0$ .

*Proof.* (1) $\Rightarrow$ (2). Let  $M^{(H)}(p) \in \langle M^{(H)}(f) | f \in I_{\mathcal{A}} \rangle$  for some  $p \in I_{\mathcal{A}}$ . Since  $G$  is an HSG-basis, by (9), there are some  $h_1, \dots, h_s \in \mathcal{A}$  so that

$$\deg\left(\sum_{j=1}^s M^{(H)}(h_{ij})M^{(H)}(g_j)\right) > \deg\left(\sum_{j=1}^s M^{(H)}(h_{i+1,j})M^{(H)}(g_j)\right), \quad i = 1, 2, \dots, d. \quad (14)$$

Hence,

$$\deg(f) = \max_i \left( \deg\left(\sum_{j=1}^s M^{(H)}(h_{ij})M^{(H)}(g_j)\right) \right). \quad (15)$$

(11) and (15) give the HSG representation.  $\square$

The second major example of gradings leads to the SAGBI-Gröbner basis concept. Here,  $\Gamma = \mathbb{N}^n$  with component-wise addition equipped with a total ordering satisfying (11). In addition,  $\gamma \geq 0, \forall \gamma \in \Gamma$ . For arbitrary  $\gamma = (\gamma_1, \dots, \gamma_n) \in \Gamma$ , the space  $P_{\gamma}^{(\Gamma)}$  is a vector space of dimension 1, namely,

$$P_{\gamma}^{(\Gamma)} = \{c \cdot x^{\gamma_1} \dots x^{\gamma_n} | c \in K\}. \quad (16)$$

The maximal part  $M^{(\Gamma)}(f)$  of a polynomial  $f$  is a product of a leading coefficient  $LC(f)$  and a leading monomial  $LM(f)$ , that is  $M^{(\Gamma)}(f) = LC(f) \cdot LM(f)$ , where  $LC(f) \in K$ . The *si*-reduction  $f \xrightarrow{G_{\mathcal{A}}} \tilde{f}$  is defined if there exists a polynomial  $g \in G$  and  $a \in \mathcal{A}$  such that  $LM(f) = LM(g)LM(a)$  and then we set

$$p = \sum_{i=1}^s h_i g_i \text{ and } M^{(H)}(p) = M^{(H)}\left(\sum_{i=1}^s h_i g_i\right) = \sum_{i \in J} M^{(H)}(h_i)M^{(H)}(g_i) \in \langle M^{(H)}(g_1), \dots, M^{(H)}(g_s) \rangle, \text{ where } J = \{i | \deg(h_i g_i) = \deg(p)\}.$$

(2) $\Rightarrow$ (3). Let  $0 \neq f \in I_{\mathcal{A}}$ . By using Algorithm 1, we get  $f \xrightarrow{F_{\mathcal{A}}} h_1 \xrightarrow{F_{\mathcal{A}}} h_2 \dots \xrightarrow{F_{\mathcal{A}}} h$ , where  $h$  is *si*-reduced any further with respect to  $F$ .  $M^{(H)}(f) \in \langle M^{(H)}(g_1), \dots, M^{(H)}(g_s) \rangle$  implies  $M^{(H)}(f) = \sum_{i \in J} M^{(H)}(h_i)M^{(H)}(g_i)$ ; then,  $f \xrightarrow{G_{\mathcal{A}}} \tilde{f} = f - \sum h_i g_i \in I_{\mathcal{A}}$ . If we follow the above process inductively, then  $f \xrightarrow{G_{\mathcal{A}}, *} 0$ .

(3) $\Rightarrow$ (1). Let

$$f_0 \xrightarrow{G_{\mathcal{A}}} f_1 \xrightarrow{G_{\mathcal{A}}} \dots \xrightarrow{G_{\mathcal{A}}} f_d = 0, \quad (11)$$

where  $M^{(H)}(f_{i-1}) = \sum_{j=1}^s M^{(H)}(h_{ij})M^{(H)}(g_j)$ ,  $i = 1, 2, \dots, d$ ,  $\deg(f_{i-1}) > \deg(f_i)$ . Then,

$$f = \sum_{j=1}^s \sum_{i=1}^d M^{(H)}(h_{ij})M^{(H)}(g_j). \quad (12)$$

Note that

$$\deg(f) = \deg(f_0) = \deg\left(\sum_{j=1}^s M^{(H)}(h_{1j})M^{(H)}(g_j)\right), \quad (13)$$

and

$\tilde{f} := (f - (M^{(\Gamma)}(f)) / (M^{(\Gamma)}(g)M^{(\Gamma)}(a))ag$ . The relation  $\xrightarrow{G_{\mathcal{A}}, *}$  is constructed as above.

A SAGBI-Gröbner basis  $G$  (with respect to a given monomial ordering and a given ideal  $I_{\mathcal{A}}$  in a subalgebra  $\mathcal{A}$ ) is a set of polynomials generating  $I_{\mathcal{A}}$  and satisfying one of the following equivalent conditions:

- (i) Every  $f \in I_{\mathcal{A}}$  has a representation:

$$f = \sum_{i=1}^s h_i g_i, \quad (17)$$

$$LM(f) = \max_{i=1}^s \{LM(h_i)LM(g_i)\},$$

where  $h_i \in \mathcal{A}$  and  $g_i \in G$ .

- (ii)  $\langle M^{(\Gamma)}(g) | g \in G \rangle = \langle M^{(\Gamma)}(f) | f \in I_{\mathcal{A}} \rangle$ .

- (iii) Every  $f \in I_{\mathcal{A}}$  *si*-reduces to 0 with respect to  $G$ .

The proof of this equivalence and many other equivalent conditions can be found in [5]. If a monomial ordering is compatible with the semiordering by degrees,

$$\deg(x^{\gamma}) > \deg(x^{\beta}) \Rightarrow \gamma > \beta, \gamma, \beta \in \mathbb{N}^n, \quad (18)$$

then any SAGBI-Gröbner representation as given in (i) is an HSG representation; in other words, a SAGBI-Gröbner basis with respect to a degree compatible ordering is an HSG-basis as well. The converse is false, as the following example shows.

*Example 1.* Let  $f_1 = x^4 + 2x^2y^2 + y^4 - 1$ ,  $f_2 = x^2y^2 + y^4 - 2$ ,  $f_3 = 2x^2 + y^2$ . These polynomials belong to the subalgebra  $\mathcal{A} = \mathbb{Q}[x^2, y^2]$ . Then, we can see that  $f_1, f_2$ , and  $f_3$  already constitute an HSG-basis for ideal  $I_{\mathcal{A}} = \langle f_1, f_2, f_3 \rangle$  in  $\mathcal{A}$ . If we order the monomials by degree lexicographical ordering, then  $\langle M^{(H)}(f) | f \in I_{\mathcal{A}} \rangle_{\mathbb{Q}[M^{(H)}(\mathcal{A})]} = \langle x^4, x^2y^2, x^2 \rangle_{\mathbb{Q}[M^{(H)}(\mathcal{A})]}$ . Every SAGBI-Gröbner basis  $G$  with respect to this ordering contains at least four elements, for instance,  $G = \{g_1, g_2, g_3, g_4\}$  with  $g_1 = x^4 + 2x^2y^2 + y^4 - 1 = f_1$ ,  $g_2 = x^2y^2 + y^4 - 2 = f_2$ ,  $g_3 = 2x^2 + y^2 = f_3$ , and  $g_4 = y^4 - 4$ . Obviously, this SAGBI-Gröbner basis is an HSG-basis as well.

### 3. Construction of HSG-Bases

In this section, we present an HSG-basis criterion, through which we can construct HSG-basis. For this purpose, we fix some notations which are necessary for this construction. Let  $\mathcal{A}$  be a  $K$ -subalgebra of  $K[x_1, \dots, x_n]$ .

- (i) We denote  $\mathcal{A} \oplus \dots \oplus \mathcal{A}$  ( $s$ -times) by  $\oplus_s \mathcal{A}$ .
- (ii) For a subset  $G \subseteq \mathcal{A}$ , we denote  $\{M^{(H)}(g_i) | g_i \in G\}$  by  $M^{(H)}(G)$ .

*Definition 3.* For  $K$ -subalgebra  $\mathcal{A}$  of  $K[x_1, \dots, x_n]$  and a subset  $G = \{g_1, \dots, g_s\} \subseteq \mathcal{A}$ ,

- (1)  $\text{syz}_{\mathcal{A}}(G) = \{\vec{a} = (a_i)_{i=1}^s \in \oplus_s \mathcal{A} | \sum_{i=1}^s a_i g_i = 0\}$ . We call an element of  $\text{syz}_{\mathcal{A}}(G)$  an  $\mathcal{A}$ -syzygy of  $G$ .
- (2) For  $\vec{a} = (a_i)_{i=1}^s \in \oplus_s \mathcal{A}$ , let  $M^{(H)}(\vec{a})$  represent the vector  $(M^{(H)}(a_i)_{i=1}^s)$ .

*Definition 4.* We call a subset  $Q = \{\vec{q}_1, \vec{q}_2, \dots, \vec{q}_m\}$  a  $M^{(H)}$ -generating set for  $\text{syz}(M^{(H)}(G))$  if  $\{M^{(H)}(\vec{q}_i) | 1 \leq i \leq m\}$  generates the  $K[M^{(H)}(\mathcal{A})]$ -module  $\text{syz}[M^{(H)}(G)]$ , i.e., for  $\vec{a} \in \text{syz}[M^{(H)}(G)]$ , there are some  $h_1, h_2, \dots, h_m \in M^{(H)}(\mathcal{A})$  such that

$$M^{(H)}(a_i)_{i=1}^s = \sum_{j=1}^m M^{(H)}(h_j) M^{(H)}(q_{ij})_{i=1}^s. \quad (19)$$

In the case of SAGBI-Gröbner bases, there is an algorithm for computing SAGBI-Gröbner bases by means of syzygies (see [6]) where syzygies and their connection to SAGBI-Gröbner bases are studied in detail. The analogue for constructing HSG-bases by means of syzygies is connected to the following result [7].

**Theorem 2 (HSG-basis criterion).** Let  $G = \{g_1, \dots, g_s\}$  be the subset of a subalgebra  $\mathcal{A}$ . Let  $Q$  be  $M^{(H)}$ -generating set for the  $\text{syz}(M^{(H)}(G))$ . Then,  $G$  is an HSG-basis for  $\langle G \rangle_{\mathcal{A}}$  if and only if for every  $\vec{q}_j = (q_{j,1}, \dots, q_{j,s}) \in Q$ , we have  $\sum_{i=1}^s q_{j,i} g_i \xrightarrow{G, \mathcal{A}} 0$ .

*Proof.*  $\Rightarrow$ : The statement is a direct result of Theorem 1.

$\Leftarrow$ : Take  $f \in \langle G \rangle_{\mathcal{A}}$ . We need to show that  $M^{(H)}(f) \in \langle M^{(H)}(G) \rangle_{K[M^{(H)}(\mathcal{A})]}$ . For this, we write  $f = \sum_{i=1}^m a_i g_i$  such that  $p_0 = \max[M^{(H)}(a_i g_i)]$  (degree wise) is minimal among all such representations of  $f$ . We have  $M^{(H)}(f) \leq p_0$ . Suppose that  $M^{(H)}(f) < p_0$ . Assume that  $a_1 g_1, \dots, a_{m_0} g_{m_0}$  are contributing to  $p_0$ , i.e.,  $M^{(H)}(a_i g_i) = p_0$  for all  $1 \leq i \leq m_0$ . If we set  $\vec{a} = (a_1, \dots, a_{m_0}, 0, \dots, 0)$ , we can see that  $M^{(H)}(\vec{a}) \in \text{syz}(M^{(H)}(G))$ . This implies that there are  $b_1, \dots, b_n \in \mathcal{A}$  and  $\vec{Q}_1, \dots, \vec{Q}_n \in \vec{Q}$  such that  $M^{(H)}(\vec{a}) = \sum_{j=1}^n M^{(H)}(b_j) M^{(H)}(\vec{Q}_j)$ . We may assume that  $M^{(H)}(b_j) M^{(H)}(q_{j,i}) M^{(H)}(g_i) = p_0$  for each  $j$  by homogeneity of the syzygies. Now,

$$\begin{aligned} f &= \sum_{i=1}^m a_i g_i - \sum_{i=1}^m \left( \sum_{j=1}^n b_j q_{j,i} \right) g_i + \sum_{j=1}^n b_j \left( \sum_{i=1}^m q_{j,i} g_i \right) \\ &= \sum_{i=1}^m \left( a_i - \sum_{j=1}^n b_j q_{j,i} \right) g_i + \sum_{j=1}^n b_j \left( \sum_{i=1}^m p_{j,i} g_i \right), \end{aligned} \quad (20)$$

where  $\sum_{i=1}^m p_{j,i} g_i$  is an HSG representation for  $\sum_{i=1}^m q_{j,i} g_i$  since  $\sum_{i=1}^m q_{j,i} g_i \xrightarrow{G} 0$ . If we define  $H_j = \max(M^{(H)}(p_{j,i} g_i))$ , then

$$H_j = M^{(H)}\left(\sum_{i=1}^m q_{j,i} g_i\right) < \max(M^{(H)}(q_{j,i} g_i)), \quad \text{for all } j, \quad (21)$$

because  $M^{(H)}(\vec{Q}_j) \in \text{syz}(M^{(H)}(G))$ .

Consider the first sum of equation (20). For  $i \leq m_0$ , we have  $M^{(H)}(a_i) = M^{(H)}\left(\sum_{j=1}^n b_j q_{j,i}\right)$ , so by the cancellation of highest terms,

$$M^{(H)}\left[\left(a_i - \sum_{j=1}^n b_j q_{j,i}\right) g_i\right] < M^{(H)}(a_i g_i) = p_0. \quad (22)$$

For  $i > m_0$ ,  $M^{(H)}(a_i g_i) < p_0$  and  $\sum_{j=1}^n M^{(H)}(b_j) M^{(H)}(q_{j,i}) = 0$  implies that

$$M^{(H)}\left(\sum_{j=1}^n b_j q_{j,i} g_i\right) < \max_j(M^{(H)}(b_j q_{j,i} g_i)) = p_0. \quad (23)$$

Since

$$M^{(H)}\left[\left(a_i - \sum_{j=1}^n b_j q_{j,i}\right) g_i\right] \leq \max\left\{M^{(H)}(a_i g_i), M^{(H)}\left(\sum_{j=1}^n b_j q_{j,i} g_i\right)\right\} < p_0 \quad (\forall i). \quad (24)$$

So, first sum of equation (20) is less than  $p_0$ . For the second sum of equation (20), we have

$$\begin{aligned} M^{(H)}\left(\sum_{j=1}^n b_j \sum_{i=1}^m p_{j,i} g_i\right) &\leq \max_{i,j} M^{(H)}(b_j p_{j,i} g_i) \\ &\leq \max_j [M^{(H)}(b_j) H_j] \\ &< \max_{i,j} (M^{(H)}(b_j q_{j,i} g_i)) = p_0. \end{aligned} \tag{25}$$

Hence, equation (20) does provide a new representation for  $f$  such that  $\max(M^{(H)}(a_i g_i)) < p_0$ , a contradiction. Therefore,  $M^{(H)}(f) = p_0$  and  $M^{(H)}(f) = \sum_{i=1}^{m_0} M^{(H)}(a_i g_i) \in \langle M^{(H)}(G) \rangle$ .  $\square$

On the basis of Theorem 2, now we present an algorithm which computes HSG-basis from a given set of generators. This algorithm is not necessarily terminating but does terminate, if and only if, the considered ideal in the subalgebra has a finite HSG-basis.

Now we present some examples which show the computation of HSG-basis through Algorithm 2.

*Example 2.* Let the subalgebra  $\mathcal{A} = Q[x^2, xy]$  and  $G = \{x^3y + x^2, xy + 2\} \subseteq \mathcal{A}$ . Consider  $H = G$ ; then,  $M^{(H)}(H) = \{x^3y, xy\}$ .

First pass through the while loop:

- (i)  $M^{(H)}(q_1)(x^3y) + M^{(H)}(q_2)(xy) = 0$  implies  $Q = \{-1, x^2\}$ . Then,  $(-1)(x^3y + x^2) + (x^2)(xy + 2) = -x^3y - x^2 + x^3y + 2x^2 = x^2$  gives  $P = \{x^2\}$ .
- (ii) As  $x^2$  is *si*-reduced with respect to  $H$ ,  $\text{red}(P) = \{x^2\}$ .
- (iii) Define:  $\text{Old}(H) = H \cup \{x^2\}$ .  
As  $H \neq \text{Old}(H)$ , we repeat the whole process. Now we have  $M^{(H)}(H) = \{x^3y, xy, x^2\}$ .

Second pass through the while loop:

- (i)  $M^{(H)}(q_1)(x^3y) + M^{(H)}(q_2)(xy) + M^{(H)}(q_3)(x^2) = 0$  implies  $(-1)(x^3y) + (0)(xy) + (xy)(x^2) = 0$ . Therefore,  $Q = \{-1, x^2, 0, (-1, 0, xy)\}$ . Then,  $(-1)(x^3y + x^2) + (0)(xy + 2) + (xy)(x^2) = -x^3y - x^2 + 0 + x^3y = -x^2$  gives  $P = \{x^2, -x^2\}$ .
- (ii) Now,  $\text{red}(P) = \emptyset$ .

Since  $\text{Old}(H) = H$ , we stop here. The HSG-basis for  $\langle G \rangle_{\mathcal{A}}$  is  $\{x^3y + x^2, xy + 2, x^2\}$ .

*Example 3.* Let  $\mathcal{A} = Q[x^2, xy]$  and  $G = \{x^3y + x^2y^2 + x^2, xy + 2\} \subseteq \mathcal{A}$ . Consider  $H = G$ ; then,  $M^{(H)}(H) = \{x^3y + x^2y^2, xy\}$ .

First pass through the while loop:

- (i)  $M^{(H)}(q_1)(x^3y + x^2y^2) + M^{(H)}(q_2)(xy) = 0$  gives  $Q = \{-1, x^2 + xy\}$ . Then, from  $(-1)(x^3y + x^2y^2 + x^2) + (x^2 + xy)(xy + 2) = -x^3y - x^2y^2 - x^2 + x^3y + x^2y^2 + 2x^2 + 2xy = x^2 + 2xy$ ,
- (ii)  $\text{red}(P) = \{x^2 - 4\}$ .
- (iii) Define:  $\text{Old}(H) = H \cup \{x^2 - 4\}$ .

As  $H \neq \text{Old}(H)$ , we repeat the whole process. Now we have  $M^{(H)}(H) = \{x^3y + x^2y^2, xy, x^2\}$ .

Second pass through the while loop:

- (i) From the equation  $M^{(H)}(q_1)(x^3y + x^2y^2) + M^{(H)}(q_2)(xy) + M^{(H)}(q_3)(x^2) = 0$ , we have  $Q = \{-1, xy, xy, (-1, x^2 + xy, 0)\}$ . We can compute  $P$  from  $(-1)(x^3y + x^2y^2 + x^2) + (xy)(xy + 2) + (xy)(x^2 - 4) = -x^3y - x^2y^2 - x^2 + x^2y^2 + 2xy + x^3y - 4xy = -x^2 - 2xy$ .
- (ii) Now,  $\text{red}(P) = \emptyset$ .

Since  $\text{Old}(H) = H$ , we stop here. The HSG-basis for  $\langle G \rangle_{\mathcal{A}}$  is  $\{x^3y + x^2y^2 + x^2, xy + 2, x^2 - 4\}$ .

## 4. Conclusion

In this paper, we presented the theory of HSG-bases, which are a good basis of an ideal in a subalgebra of a polynomial ring. We can further develop this theory for an arbitrary grading for which HSG-bases would be a special case for degree-based grading.

## Data Availability

No data are required to support the study.

## Conflicts of Interest

The authors declare that they have no conflicts of interest.

## References

- [1] F. S. Macaulay, "The algebraic theory of modular systems," *Cambridge Tracts in Mathematics and Mathematical Physics*, Cambridge University Press, Cambridge, UK, 1916.
- [2] T. Sauer, "Gröbner bases, H-bases and interpolation," *Transactions of the American Mathematical Society*, vol. 353, no. 6, pp. 2293–2308, 2000.
- [3] H. M. Möller and T. Sauer, "H-bases for polynomial interpolation and system solving," *Advances in Computational Mathematics*, vol. 12, no. 4, pp. 335–362, 2000.
- [4] J. A. Khan, M. A. Binyamin, and S. Rabia, "Subalgebra analogue to H-basis for ideals," *Hacettepe Journal of Mathematics and Statistics*, vol. 6, no. 45, pp. 1685–1692, 2016.
- [5] L. Robbiano and M. Sweedler, *Subalgebra Bases in Commutative Algebra*, Springer, Berlin, Germany, 1990.
- [6] J. L. Miller, *Effective Algorithms for Intrinsically Computing SAGBI-Gröbner Bases in a Polynomial Ring over a Field*, pp. 421–433, Cambridge University Press, Cambridge, UK, 1998.
- [7] J. L. Miller, "Analogues of gröbner bases in polynomial rings over a ring," *Journal of Symbolic Computation*, pp. 139–153, 2008.